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The Recursive Structures of Manin Symbols over \mathbb{Q} , Cusps and Elliptic Points on $X_0(N)$

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Abstract: Firstly, we present a more explicit formulation of the complete system $D(N)$ of representatives of Manin’s symbols over \mathbb{Q} , which was initially given by Shimura. Then, we establish a bijection between $D(M) \times D(N)$ and $D(MN)$ for $(M, N) = 1$, which reveals a recursive structure between Manin’s symbols of different levels. Based on Manin’s complete system $\Pi(N)$ of representatives of cusps on $X_0(N)$ and Cremona’s characterization of the equivalence between cusps, we establish a bijection between a subset $C(N)$ of $D(N)$ and $\Pi(N)$, and then establish a bijection between $C(M) \times C(N)$ and $C(MN)$ for $(M, N) = 1$. We also provide a recursive structure for elliptic points on $X_0(N)$. Based on these recursive structures, we obtain recursive algorithms for constructing Manin symbols over \mathbb{Q} , cusps, and elliptic points on $X_0(N)$. This may give rise to more efficient algorithms for modular elliptic curves. As direct corollaries of these recursive structures, we present a recursive version of the genus formula and prove constructively formulas of the numbers of $D(N)$, cusps, and elliptic points on $X_0(N)$.

Keywords: modular curve; elliptic curve; recursive structure; Manin’s symbols over \mathbb{Q} ; cusps; elliptic points; algorithmic number theory

MSC: 11A05; 11F06; 20H05; 20J05



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1. Introduction

In his seminal monograph [1] (Chapter 1, Proposition 1.43), G. Shimura defined a complete set $D(N)$ of representatives for the projective line $\mathbb{P}^1(\mathbb{Z}/N\mathbb{Z})$ over $\mathbb{Z}/N\mathbb{Z}$ to be all couples $\{c, d\}$ of positive integers satisfying

$$(*) \quad (c, d) = 1, d|N, 1 \leq c \leq N/d \text{ (or } c \text{ in any set of representatives for } \mathbb{Z} \text{ modulo } (N/d)),$$

where (c, d) denote the greatest common divisor of integers c and d .

Let $[x]$ be the greatest integer less than or equal to x . For two integers a, b with $b \neq 0$, define

$$\left[\frac{a}{b}\right]' = \begin{cases} \frac{a}{b} - 1 & \text{if } b|a, \\ \left[\frac{a}{b}\right] & \text{otherwise,} \end{cases}$$

then $1 \leq a - b\left[\frac{a}{b}\right]' \leq b$. In this paper, we define

$$D(N) = \{(c, d) : c, d \in \mathbb{Z}, c, d \geq 1, c|N, (c, d) = 1 \text{ and } (c, d - \frac{N}{c}(\left[\frac{cd}{N}\right]' - n)) \geq 2 \text{ for } 0 \leq n < \left[\frac{cd}{N}\right]'\}. \quad (1)$$

We then establish a bijection between $D(M) \times D(N)$ and $D(MN)$ for $(M, N) = 1$ in Section 2. This result gives a recursive algorithm to construct the projective line $\mathbb{P}^1(\mathbb{Z}/N\mathbb{Z})$ over $\mathbb{Z}/N\mathbb{Z}$.

Let $\Pi(N) = \{[\delta; a \bmod (\delta, N\delta^{-1})] : a, \delta \in \mathbb{Z}, \delta \geq 1, \delta|N, 1 \leq a \leq (\delta, N\delta^{-1})\}$. In [2] (Proposition 2.2), Manin proved that there exists a bijection between $\Pi(N)$ and the set of cusps on $X_0(N)$. Based on Manin’s result and Cremona’s characterization (See Proposition 3), we identify $\Pi(N)$ with

$$C(N) = \{(c, d) : c, d \in \mathbb{Z}, 1 \leq c \leq N, c|N, (c, d) = 1 \text{ and } (c, d - (c, Nc^{-1})[\frac{d}{(c, Nc^{-1})}]' + \frac{Nn}{c}) \geq 2 \text{ for } 0 \leq n < \frac{c(c, Nc^{-1})}{N}[\frac{d}{(c, Nc^{-1})}]'\}, \tag{2}$$

which is a subset of $D(N)$. In Section 3, we establish a bijection between $C(N_1N_2)$ and $C(N_1) \times C(N_2)$ for $(N_1, N_2) = 1$. This result gives a recursive algorithm to construct the complete set of representatives of $\Gamma_0(N)$ -inequivalent cusps.

Define

$$\begin{aligned} E_2(N) &= \{(1, d) : (1, d) \in D(N), 1 + d^2 \equiv 0 \pmod{N}\}, \\ E_3(N) &= \{(1, d) : (1, d) \in D(N), 1 - d + d^2 \equiv 0 \pmod{N}\}. \end{aligned} \tag{3}$$

Then, there exist bijections between $E_2(N), E_3(N)$ and complete sets of representatives of $\Gamma_0(N)$ -inequivalent elliptic points of order 2 and 3, respectively. In Section 4, we establish bijections between $E_2(N_1N_2)$ and $E_2(N_1) \times E_2(N_2), E_3(N_1N_2)$ and $E_3(N_1) \times E_3(N_2)$, for $(N_1, N_2) = 1$. These results give a recursive algorithm for constructing the complete set $E_3(N)$ and $E_2(N)$ of $\Gamma_0(N)$ -inequivalent elliptic points of order 2, 3.

The elements in $\mathbb{P}^1(\mathbb{Z}/N\mathbb{Z})$ are called Manin symbols [3] (Section 2.2) and there exists a bijection between the set of right cosets of $\Gamma_0(N)$ in $SL(2, \mathbb{Z})$ and $\mathbb{P}^1(\mathbb{Z}/N\mathbb{Z})$ [2] (Proposition 2.4). An important step in the modular elliptic algorithm is to construct a complete set of representatives for the projective line $\mathbb{P}^1(\mathbb{Z}/N\mathbb{Z})$ and a complete set of representatives of $\Gamma_0(N)$ -inequivalent cusps [3] (Chapter II). The recursive structure of $D(N), C(N), E_2(N)$ and $E_3(N)$ may give rise to a more efficient modular elliptic algorithm.

As direct corollaries of these recursive structures, we present a recursive version of the genus formula and elementary proofs of formulas of the numbers $\mu(N), v_\infty(N), v_2(N)$ and $v_3(N)$ of $D(N), C(N), E_2(N), E_3(N)$. Note that Schoeneberg’s proof and Shimura’s proof for formulas of $\mu(N), v_\infty(N), v_2(N)$ and $v_3(N)$ use the theory of quadratic fields, see [4] (Chapter IV, Section 8) and [1] (Chapter 1, Proposition 1.43). Their proofs may make these formulas hard to approach when compared with our proofs.

2. The Recursive Structure of Manin Symbols over \mathbb{Q}

We firstly give some necessary notations and facts, for details, see [3].

Definition 1.

- (a) $D_2(N) = \{(c, d) : c, d \in \mathbb{Z}, (c, d, N) = 1\}$;
- (b) $\forall (c_1, d_1), (c_2, d_2) \in D_2(N)$, define $(c_1, d_1) \sim (c_2, d_2)$ if $c_1d_2 \equiv d_1c_2 \pmod{N}$, then \sim is an equivalence relation on $D_2(N)$;
- (c) $\forall (c, d) \in D_2(N)$, define $(c : d) = \{(c', d') : (c', d') \in D_2(N), (c', d') \sim (c, d)\}$;
- (d) $\mathcal{D}(N) = D_2(N) / \sim = \{(c : d) : (c, d) \in D_0(N)\}$;
- (e) $D_1(N) = \{(c, d) : c, d \in \mathbb{Z}, c, d \geq 1, c|N, (c, d, \frac{N}{c}) = 1, cd \leq N\}$;
- (f) $D(N)$ is defined in (1);
- (g) $\mu(N), v_\infty(N), v_2(N)$ and $v_3(N)$ are the numbers of elements in $D(N), C(N), E_2(N)$ and $E_3(N)$, respectively.

As pointed out by a referee, the index $\mu(N)$ of $\Gamma_0(N)$ in $SL(2, \mathbb{Z})$ is called the Dedekind psi function, usually denoted $\psi(N)$, see [5,6]. Here, we follow Shimura’s notations in [1] (Proposition 1.43).

Lemma 1. *Let $c, d, h \in \mathbb{Z}$, $(c, d, h) = 1, c, d \geq 1$ and $d \leq h$, then there exists an integer k such that $(c, d + hk) = 1$ and $0 \leq k < c$.*

Proof. If $c = 1$, take $k = 0$ then $(c, d + hk) = 1$. Thus, let $c \geq 2$ in the following. Let $c = p_1^{\alpha_1} \cdots p_s^{\alpha_s}$ be the standard factorization of c . The proof is by induction on the numbers of distinct prime divisors in c . Suppose that $c = p_1^{\alpha_1}$. Assume that $(p_1^{\alpha_1}, d) \geq 2$ and $(p_1^{\alpha_1}, d + h) \geq 2$ then $p_1|d$ and $p_1|(d + h)$. Thus, $p_1|d$ and $p_1|h$, this contradicts with $(c, d, h) = 1$, and hence $(c, d + hk) = 1$ for some $0 \leq k \leq 1 < c$.

Let $c_1 = p_1^{\alpha_1} \cdots p_{s-1}^{\alpha_{s-1}}$. By the induction hypothesis, there exists an integer k_1 such that $(c_1, d + hk_1) = 1$ and $0 \leq k_1 < c_1$. Then, $(c_1, d + hk_1 + hc_1) = 1$. Assume that $(p_s^{\alpha_s}, d + hk_1) \geq 2$ and $(p_s^{\alpha_s}, d + hk_1 + hc_1) \geq 2$ then $p_s|(d + hk_1)$ and $p_s|(d + hk_1 + hc_1)$. Thus, $p_s|hc_1$ and hence $p_s|h$ by $(p_s, c_1) = 1$. Therefore, $p_s|d$. This contradicts with $(c, d, h) = 1$ and hence $(c, d + hk_1) = 1$ or $(c, d + hk_1 + hc_1) = 1$. Take $k = k_1$ or $k = c_1 + k_1$, then $(c, d + hk) = 1$ for some $0 \leq k_1 \leq k \leq c_1 + k_1 < 2c_1 \leq c$. This completes the proof by the induction principle. \square

Corollary 1. *Let $a, b, c \in \mathbb{Z}$, $(a, b, c) = 1$, then the equation $ax + by + cyz = 1$ has solutions in \mathbb{Z} .*

Lemma 2. *There exists a bijection between $D(N)$ and $D_1(N)$.*

Proof. Let $(c, d) \in D(N)$. Define $d_n = d - \frac{N}{c}[\frac{cd}{N}]' + \frac{Nn}{c}$ for all $n \in \mathbb{Z}$. Then, $1 \leq d_0 \leq \frac{N}{c}$ and $(c, d_0, \frac{N}{c}) = 1$ by $(c, d) = 1$. Thus $(c, d_0) \in D_1(N)$. Define $\Phi : D(N) \rightarrow D_1(N)$ by sending (c, d) to (c, d_0) .

Let $(u, v) \in D(N)$ such that $\Phi(c, d) = \Phi(u, v)$. Define $v_n = v - \frac{N}{u}[\frac{uv}{N}]' + \frac{Nn}{u}$ for all $n \in \mathbb{Z}$. Then, $c = u$ and $d_0 = v_0$. Thus, $d_n = v_n$ for all $n \in \mathbb{Z}$. Let $e = [\frac{cd}{N}]'$ and $w = [\frac{cv}{N}]'$. Then, $d = d_e$ and $v = v_w$. Suppose that $e < w$ then $(c, d_e) = 1$ by $(c, d) = 1$ but $(c, d_e) \geq 2$ by $(c, v) \in D(N)$ and $d_e = v_e$, a contradiction and thus $e \geq w$. $e \leq w$ holds by a similar proof and thus $e = w$ and $(c, d) = (u, v)$. Therefore, Φ is an injection from $D(N)$ to $D_1(N)$.

Let $(c, d_0) \in D_1(N)$. By Lemma 1, there exists an integer k such that $(c, d_0 + \frac{Nk}{c}) = 1$ and $0 \leq k \leq c - 1$. Let $0 \leq k_0 \leq k$ such that $(c, d_0 + \frac{Nk_0}{c}) = 1$ and $(c, d_0 + \frac{Nn}{c}) = 1$ for all $0 \leq n < k_0$. Define $d = d_0 + \frac{Nk_0}{c}$. Then, $(c, d) \in D(N)$ and $\Phi((c, d)) = (c, d_0)$. Therefore, Φ is a surjection from $D(N)$ to $D_1(N)$. \square

Lemma 3. *There exists a bijection between $\mathcal{D}(N)$ and $D(N)$, i.e., $D(N)$ is a complete system of the representatives of elements of $\mathcal{D}(N)$.*

Proof. Define $\Phi : D(N) \rightarrow \mathcal{D}(N)$ by the natural map, i.e., $\Phi((c, d)) = (c : d)$.

Let $(c : d) \in \mathcal{D}(N)$. Then, $(c, d, N) = 1$. Define $c_1 = (c, N)$, d_0 to be the unique solution of the congruence equation $\frac{c}{c_1}x \equiv d \pmod{\frac{N}{c_1}}$ such that $1 \leq d_0 \leq \frac{N}{c_1}$. Then, there exists an integer y such that $\frac{c}{c_1}d_0 + \frac{N}{c_1}y = d$. Assume that there exists a prime p such

that $p|(c_1, d_0, \frac{N}{c_1})$. Then, $p|d$ and $p|(c, N)$, this contradicts with $(c, d, N) = 1$, and thus $(c_1, d_0, \frac{N}{c_1}) = 1$. Hence, $(c_1, d_0) \in D_1(N)$. Then, there exists the unique $(c_1, d_1) \in D(N)$ which corresponds to (c_1, d_0) . Hence, $(c_1, d_1) \in (c : d)$, i.e., $\Phi((c_1, d_1)) = (c : d)$.

Assume that $(c_1, d_1), (c_2, d_2) \in D(N)$ such that $\Phi((c_1, d_1)) = \Phi((c_2, d_2))$. Then, $(c_1 : d_1) = (c_2 : d_2)$ and thus there exists an integer k such that $c_1 d_2 - c_2 d_1 = Nk$. Thus, $c_1|c_2 d_1$ by $c_1|N$ and $c_2|c_1 d_2$ by $c_2|N$. Hence, $c_1|c_2$ by $(c_1, d_1) = 1$ and $c_2|c_1$ by $(c_2, d_2) = 1$. Therefore, $c_1 = c_2$ and $d_1 = d_2$ by $d_1 \equiv d_2 \pmod{\frac{N}{c_1}}$ and the definition of $D(N)$. Thus, Φ is a bijection between $\mathcal{D}(N)$ and $D(N)$. This completes the proof. \square

Theorem 1. Let $M, N \in \mathbb{Z}, M, N \geq 1, (M, N) = 1$. Then, there exists a bijection between $D(M) \times D(N)$ and $D(MN)$.

Proof. Let $(a, b) \in D(M)$ and $(c, d) \in D(N)$. Assume that there exists a prime p such that $p|(ac, bN + dM - \frac{MN}{ac}[\frac{ac(bN + dM)}{MN}]', \frac{MN}{ac})$. Then, $p|ac, p|\frac{MN}{a c}$ and

$$p|bN + dM - \frac{MN}{ac}[\frac{ac(bN + dM)}{MN}]'.$$

Then $p|a, p|\frac{M}{a}$ or $p|c, p|\frac{N}{c}$ by $(M, N) = 1, a|M, c|N$. If $p|a, p|\frac{M}{a}$ then $p|bN$ and thus $p|N$ by $(a, b) = 1$, which contradicts with $(M, N) = 1$. The case of $p|c, p|\frac{N}{c}$ is tackled in a similar way. Therefore $(ac, bN + dM - \frac{MN}{ac}[\frac{ac(bN + dM)}{MN}]', \frac{MN}{ac}) = 1$ and

$$(ac, bN + dM - \frac{MN}{ac}[\frac{ac(bN + dM)}{MN}]') \in D_1(MN).$$

Define $e = ac, f = bN + dM - \frac{MN}{ac}([\frac{ac(bN + dM)}{MN}]' - k)$ for some k such that

$$(ac, bN + dM - \frac{MN}{ac}([\frac{ac(bN + dM)}{MN}]' - n)) \geq 2$$

for all $0 \leq n < k$. Then $(e, f) \in D(MN)$. Define $\Phi : D(M) \times D(N) \rightarrow D(MN)$ by sending $((a, b), (c, d))$ to (e, f) .

Assume that $\Phi((a, b), (c, d)) = \Phi((a_1, b_1), (c_1, d_1))$ for some $(a, b), (a_1, b_1) \in D(M)$ and $(c, d), (c_1, d_1) \in D(N)$. Then

$$\begin{aligned} & (ac, bN + dM - \frac{MN}{ac}([\frac{ac(bN + dM)}{MN}]' - k)) \\ &= (a_1 c_1, b_1 N + d_1 M - \frac{MN}{a_1 c_1}([\frac{a_1 c_1 (b_1 N + d_1 M)}{MN}]' - k_1)). \end{aligned}$$

Thus, $ac = a_1 c_1$ and

$$\begin{aligned} & bN + dM - \frac{MN}{ac}([\frac{ac(bN + dM)}{MN}]' - k) \\ &= b_1 N + d_1 M - \frac{MN}{a_1 c_1}([\frac{a_1 c_1 (b_1 N + d_1 M)}{MN}]' - k_1). \end{aligned}$$

Hence, $a = a_1, c = c_1$ by $(M, N) = 1, a|M, a_1|M, c|N, c_1|N$. Therefore,

$$\begin{aligned} & bN + dM - \frac{MN}{ac} \left(\left[\frac{ac(bN + dM)}{MN} \right]' - k \right) \\ &= b_1N + d_1M - \frac{MN}{ac} \left(\left[\frac{ac(b_1N + d_1M)}{MN} \right]' - k_1 \right). \end{aligned}$$

Thus $d \equiv d_1 \pmod{\frac{N}{c}}$ and $b \equiv b_1 \pmod{\frac{M}{a}}$ by $(M, N) = 1$. Hence $b = b_1, d = d_1$. Then $((a, b), (c, d)) = ((a_1, b_1), (c_1, d_1))$.

Let $(e, f) \in D(MN)$. Then $e|MN, (e, f) = 1$ and $(e, f - \frac{MN}{e} \left(\left[\frac{ef}{MN} \right]' - n \right)) \geq 2$ for $0 \leq n < \left[\frac{ef}{MN} \right]'$. Let $a = (e, M), c = (e, N)$, then $e = ac, a|M$ and $c|N$. Let x_0, y_0, z_0 be a particular solution of the equation

$$Nx + My + \frac{MN}{ac}z = f \tag{4}$$

then $x = \frac{M}{a}X + x_0, y = \frac{N}{c}Y + y_0, z = -cX - aY + z_0$ are solutions of (4) for all integers X, Y . Take $b_1 = x_0 - \frac{M}{a} \left[\frac{ax_0}{M} \right]', d_1 = y_0 - \frac{N}{c} \left[\frac{cy_0}{N} \right]'$, then

$$Nb_1 + Md_1 + \frac{MN}{ac} \left(c \left[\frac{ax_0}{M} \right]' + a \left[\frac{cy_0}{N} \right]' + z_0 \right) = f, 1 \leq b_1 \leq \frac{M}{a}, 1 \leq d_1 \leq \frac{N}{c}.$$

Then, $(a, b_1, \frac{M}{a}) = 1$ by $a|M, (e, f) = 1$ and $(c, d_1, \frac{N}{c}) = 1$ by $c|N, (e, f) = 1$. Hence, $(a, b_1) \in D_1(M), (c, d_1) \in D_1(N)$. Let $(a, b) \in D(M)$ and $(c, d) \in D(N)$ which correspond to (a, b_1) and (c, d_1) , respectively. Then $b = b_1 + \frac{M}{a}k_1$ and $d = d_1 + \frac{N}{c}k_2$ for some k_1, k_2 . Then $Nb + Md + \frac{MN}{ac} \left(c \left[\frac{ax_0}{M} \right]' + a \left[\frac{cy_0}{N} \right]' - ck_1 - ak_2 + z_0 \right) = f$. Then $(e, f) = \Phi((a, b), (c, d))$.

Thus, Φ is a bijection between $D(M) \times D(N)$ and $D(MN)$. \square

Proposition 1. Let p be a prime and l a positive integer. Then

- (a) $D(p^l) = \{(1, d) : 1 \leq d \leq p^l\} \cup \{(p^l, 1)\} \cup \{(p^\alpha, kp + d) : 1 \leq \alpha \leq l - 1, 1 \leq d \leq p - 1, 0 \leq k \leq p^{l-\alpha-1} - 1\}$;
- (b) $\mu(p^l) = p^l \left(1 + \frac{1}{p} \right)$;
- (c) $\mu(N) = N \prod_{p|N} \left(1 + \left(\frac{1}{p} \right) \right)$.

Proof. (c) is immediately from (b) and Theorem 1. \square

$D(MN)$ can be constructed using Algorithm 1.

Algorithm 1: $D(MN)$

- (1) Construct $D(p^l)$ by Proposition 1(a);
 - (2) Given $D(M)$ and $D(N)$ for $(M, N) = 1$, $D(MN)$ is constructed as follows. For all $(a, b) \in D(M), (c, d) \in D(N)$, define $e = ac$,
 $f = bN + dM - \frac{MN}{ac}([\frac{ac(bN + dM)}{MN}]' - k)$ for some $k \in \mathbb{Z}$ such that $(e, f) = 1$
 and $(ac, bN + dM - \frac{MN}{ac}([\frac{ac(bN + dM)}{MN}]' - n)) \geq 2$ for all $0 \leq n < k$. Then, $(e, f) \in D(MN)$ and all elements in $D(MN)$ are constructed if all pairs in $D(M) \times D(N)$ are processed.
-

3. The Recursive Structure of Cusps

In order to describe the cusps on $X_0(N)$, Ju. I. Manin in [2] introduced the set $\Pi(N)$, which consists of pairs of the form $[\delta; a \pmod{(\delta, N\delta^{-1})}]$. Here, δ runs through all positive divisors of N , and the second coordinate of the pair runs through any invertible class of residues modulo the greatest common divisor of δ and $N\delta^{-1}$. If $(\delta, N\delta^{-1}) = 1$ we sometimes put simply 1 in place of the second coordinate.

Proposition 2. Let $\delta|N, u, v \in \mathbb{Z}; (u, v\delta) = (v, N\delta^{-1}) = 1$. The map $\mathbb{Q} \cup \{i\infty\} \rightarrow \Pi(N)$ of the form $\frac{u}{v\delta} \mapsto [\delta; uv \pmod{(\delta, N\delta^{-1})}]$ gives an isomorphism of the set of cusps on $X_0(N)$ with $\Pi(N)$.

Proof. See Proposition 2.2 in [2]. \square

In [3], (Proposition 2.2.3), J. E. Cremona gives the following characterization of cusps of $X_0(N)$.

Proposition 3. For $j = 1, 2$ let $\alpha_j = p_j/q_j$ be cusps written in the lowest terms. The following are equivalent:

- (a) $\alpha_2 = M(\alpha_1)$ for some $M \in \Gamma_0(N)$;
- (b) $q_2 \equiv uq_1 \pmod{N}$ and $up_2 \equiv p_1 \pmod{(q_1, N)}$, with $(u, N) = 1$;
- (c) $s_1q_2 \equiv s_2q_1 \pmod{(q_1q_2, N)}$, where s_j satisfies $p_j s_j \equiv 1 \pmod{q_j}$.

Definition 2.

- (a) $C_1(N) = \{(c, d) : c, d \in \mathbb{Z}, 1 \leq c \leq N, c|N, 1 \leq d \leq (c, Nc^{-1}), (c, d, Nc^{-1}) = 1\}$,
- (b) $C(N)$ is defined in (2).

Lemma 4. There exists a bijection between $C_1(N)$ and $C(N)$.

Proof. It holds by $C_1(N) \subseteq D_1(N), C(N) \subseteq D(N)$ and Lemma 2. \square

Lemma 5. There exists a bijection between $\Gamma_0(N) \backslash \mathbb{Q} \cup \{i\infty\}$ and $C_1(N)$.

Proof. Let $\gamma_i = \begin{pmatrix} a_i & b_i \\ c_i & d_i \end{pmatrix}, \gamma_j = \begin{pmatrix} a_j & b_j \\ c_j & d_j \end{pmatrix} \in \text{SL}_2(\mathbb{Z})$ such that $(c_i, d_i), (c_j, d_j) \in D(N)$ for $1 \leq i < j \leq \mu(N)$ then $\text{SL}_2(\mathbb{Z}) = \Gamma_0(N)\gamma_1 \cup \dots \cup \Gamma_0(N)\gamma_{\mu(N)}$ and $\Gamma_0(N)\gamma_i \neq \Gamma_0(N)\gamma_j$. $\forall c, a \in \mathbb{Z}, (c, a) = 1, c \geq 1$, let $\begin{pmatrix} a & b \\ c & d \end{pmatrix} \in \text{SL}_2(\mathbb{Z})$ for some a, b . Then there exists $\gamma \in \Gamma_0(N), 1 \leq i \leq \mu(N)$ such that $\begin{pmatrix} a & b \\ c & d \end{pmatrix} = \gamma\gamma_i$. Thus, $\begin{pmatrix} a & b \\ c & d \end{pmatrix}(\infty) = \gamma\gamma_i(\infty)$,

$\gamma\left(\frac{a_i}{c_i}\right) = \frac{a}{c}$ and $\Gamma_0(N)\frac{a}{c} = \Gamma_0(N)\frac{a_i}{c_i}$. Then, $\Gamma_0(N)\backslash\mathbb{Q} \cup \{i\infty\} = \{\Gamma_0(N)\frac{a_i}{c_i} : 1 \leq i \leq \mu(N)\}$. Define $\Phi : \Gamma_0(N)\backslash\mathbb{Q} \cup \{i\infty\} \rightarrow C_1(N)$ by

$$\Gamma_0(N)\frac{a}{c} \mapsto (c_i, d_i - (c_i, c_i^{-1}N) \left[d_i (c_i, c_i^{-1}N)^{-1} \right]'), \Gamma_0(N) \cdot i\infty \mapsto (N, 1).$$

By Proposition 3, $\Gamma_0(N)\frac{a_i}{c_i} = \Gamma_0(N)\frac{a_j}{c_j}$ if $c_i d_j \equiv c_j d_i \pmod{(c_i c_j, N)}$. Then, $c_i d_j = c_j d_i + (c_i c_j, N)h$ for some $h \in \mathbb{Z}$. Thus, $c_i = c_j$ by $c_i | N, c_j | N, (c_i, d_i) = 1$ and $(c_j, d_j) = 1$. Hence, $c_i d_j \equiv c_j d_i \pmod{(c_i c_j, N)}$ if $d_i \equiv d_j \pmod{(c_i, c_i^{-1}N)}$. Therefore, Φ is a bijection between $\Gamma_0(N)\backslash\mathbb{Q} \cup \{i\infty\}$ and $C_1(N)$. \square

Lemma 6. *There exists a bijection between $\Gamma_0(N)\backslash\mathbb{Q} \cup \{i\infty\}$ and $C(N)$.*

Proof. It is immediately from Lemmas 4 and 5. \square

Lemma 7. *Let $(N_1, N_2) = 1$. Then, there exists a bijection between $C_1(N_1 N_2)$ and $C_1(N_1) \times C_1(N_2)$.*

Proof. Let $(c, d) \in C_1(N_1 N_2)$ then $c | N_1 N_2, d \leq (c, N_1 N_2 c^{-1}), (d, c, N_1 N_2 c^{-1}) = 1$. Let $c_1 = (c, N_1), c_2 = (c, N_2)$ then $c = c_1 c_2, (c_1, c_2) = 1$ and $(d, c_1 c_2, N_1 c_1^{-1} N_2 c_2^{-1}) = 1$. Thus, $(d, (c_1, N_1 c_1^{-1})) = 1, (d, (c_2, N_2 c_2^{-1})) = 1$ by $(c, N_1 N_2 c^{-1}) = (c_1, N_1 c_1^{-1})(c_2, N_2 c_2^{-1})$. Let $d_1 = d - (c_1, N_1 c_1^{-1})[d(c_1, N_1 c_1^{-1})^{-1}]'$ and $d_2 = d - (c_2, N_2 c_2^{-1})[d(c_2, N_2 c_2^{-1})^{-1}]'$ then $(d_1, (c_1, N_1 c_1^{-1})) = 1$ and $(d_2, (c_2, N_2 c_2^{-1})) = 1$. Thus, $(c_1, d_1) \in C_1(N_1)$ and $(c_2, d_2) \in C_1(N_2)$. Define $\Phi : C_1(N_1 N_2) \rightarrow C_1(N_1) \times C_1(N_2)$ by $(c, d) \mapsto ((c_1, d_1), (c_2, d_2))$.

For any $((c_1, d_1), (c_2, d_2)) \in C_1(N_1) \times C_1(N_2)$, let $c = c_1 c_2$ there exists an integer d such that $d \equiv d_1 \pmod{(c_1, N_1 c_1^{-1})}, d \equiv d_2 \pmod{(c_2, N_2 c_2^{-1})}$ and

$$1 \leq d \leq (c_1, N_1 c_1^{-1})(c_2, N_2 c_2^{-1}) = (c, N_1 N_2 c^{-1})$$

by $((c_1, N_1 c_1^{-1}), (c_2, N_2 c_2^{-1})) = 1$. Thus $(c, d) \in C_1(N_1 N_2)$ and hence Φ is a surjective map.

Let $\Phi((c, d)) = \Phi((c', d'))$. Then, $((c_1, d_1), (c_2, d_2)) = ((c'_1, d'_1), (c'_2, d'_2)), (c_1, d_1) = (c'_1, d'_1)$ and $(c_2, d_2) = (c'_2, d'_2)$. Thus, $c_1 = c'_1, c_2 = c'_2, d_1 = d'_1$ and $d_2 = d'_2$. Hence, $c = c_1 c_2 = c'_1 c'_2 = c'$ and $d = d'$ by $d \equiv d_1 \pmod{(c_1, N_1 c_1^{-1})}, d \equiv d_2 \pmod{(c_2, N_2 c_2^{-1})}, d' \equiv d'_1 \pmod{(c_1, N_1 c_1^{-1})}$ and $d' \equiv d'_2 \pmod{(c_2, N_2 c_2^{-1})}$. Therefore, Φ is an injective map. Then Φ is a bijection between $C_1(N_1 N_2)$ and $C_1(N_1) \times C_1(N_2)$. \square

Theorem 2. *Let $(N_1, N_2) = 1$. Then, there exists a bijection between $C(N_1 N_2)$ and $C(N_1) \times C(N_2)$.*

Proof. It is immediately from Lemmas 4 and 7. \square

Proposition 4. *Let p be a prime and l a positive integer. Then,*

- (a) $C(p^l) = \{(1, 1), (p^l, 1)\} \cup \{(p^\alpha, kp + d) : 1 \leq \alpha \leq l - 1, 1 \leq d \leq p - 1, 0 \leq k \leq p^{\min\{\alpha, l - \alpha\} - 1} - 1\}$;
- (b) $v_\infty(p^l) = \begin{cases} (p + 1)p^{\frac{l}{2} - 1} & \text{if } 2 | l, \\ 2p^{\frac{l-1}{2}} & \text{otherwise;} \end{cases}$
- (c) $v_\infty(N) = \prod_{p | N} v_\infty(p^l)$.

Proof. (c) is immediately from (b) and Theorem 2. \square

$C(N)$ can be constructed using Algorithm 2.

Algorithm 2: $C(N)$

- (1) Construct $C(p^l)$ by Proposition 4(a);
- (2) Let $N = N_1N_2$ for $(N_1, N_2) = 1$. Given $C(N_1)$ and $C(N_2)$. $C(N)$ is constructed as follows. For all $(c_1, d_1) \in C(N_1), (c_2, d_2) \in C(N_2)$, define $c = c_1c_2$. Determinate d_0 such that $d_0 \equiv d_1 \pmod{(c_1, N_1c_1^{-1})}, d_0 \equiv d_2 \pmod{(c_2, N_2c_2^{-1})}$ and

$$1 \leq d_0 \leq (c_1, N_1c_1^{-1})(c_2, N_2c_2^{-1}).$$

Determinate $d = d_0 + \frac{Nk}{c}$ such that $(c, d) = 1$ and

$(c, d_0 + \frac{Nn}{c}) \geq 2$ for $0 \leq n < k$. Then, $(c, d) \in C(N_1N_2)$ and all elements in $C(N_1N_2)$ are constructed if all pairs in $C(N_1) \times C(N_2)$ are processed.

4. The Recursive Structure of Elliptic Points of $X_0(N)$

Let $\rho = \frac{-1 + \sqrt{3}i}{2}$. $E_2(N)$ and $E_3(N)$ are defined in (3). Then,

$$\left\{ \frac{-d+i}{1+d^2} : (1, d) \in E_2(N) \right\} \text{ and } \left\{ \frac{1-2d+\sqrt{3}i}{2(1-d+d^2)} : (1, d) \in E_3(N) \right\}$$

are complete sets of representatives of $\Gamma_0(N)$ -inequivalent elliptic points of order 2, 3, respectively.

Theorem 3. Let $N_1, N_2 \in \mathbb{Z}, N_1, N_2 \geq 1$ and $(N_1, N_2) = 1$. Then

- (a) there exists a bijection between $E_3(N_1) \times E_3(N_2)$ and $E_3(N_1N_2)$;
- (b) there exists a bijection between $E_2(N_1) \times E_2(N_2)$ and $E_2(N_1N_2)$.

Proof. (a) Let $(1, d_1) \in E_3(N_1)$ and $(1, d_2) \in E_3(N_2)$. Let d be the unique integer such that $d \equiv d_1 \pmod{N_1}, d \equiv d_2 \pmod{N_2}$ and $1 \leq d \leq N_1N_2$ then $d^2 - d + 1 \equiv 0 \pmod{N_1N_2}$.

Hence, $(1, d) \in E_3(N_1N_2)$. Define

$$\Phi : E_3(N_1) \times E_3(N_2) \rightarrow E_3(N_1N_2), ((1, d_1), (1, d_2)) \mapsto (1, d).$$

Then, Φ is a bijection between $E_3(N_1) \times E_3(N_2)$ and $E_3(N_1N_2)$. The proof of (b) is similar to that of (a) and omitted. \square

Proposition 5. Let $p \in \mathbb{Z}$ be a prime and $l \in \mathbb{Z}, l \geq 1$. Then

$$v_2(p^l) = \begin{cases} 0 & \text{if } p \equiv 3 \pmod{4} \text{ or } 4|p^l, \\ 1 & \text{if } p = 2, \\ 2 & \text{if } p \equiv 1 \pmod{4}. \end{cases}$$

Proof. Let $(1, d) \in E_2(p^l)$ then $d^2 + 1 \equiv 0 \pmod{p^l}$. Since the system of two equations $x^2 + 1 \equiv 0 \pmod{p}$ and $2x \equiv 0 \pmod{p}$ has a common solution if $p = 2$, the number of solutions of $x^2 + 1 \equiv 0 \pmod{p^l}$ is equal to that of $x^2 + 1 \equiv 0 \pmod{p}$ if $p \neq 2$. The cases of $p = 2$ or $4|p^l$ are trivial and we then let $p \geq 3$ in the following. Then, $x^2 + 1 \equiv 0 \pmod{p}$

has a solution if $\left(\frac{-1}{p}\right) = 1$ if $p \equiv 1 \pmod{4}$ by $\left(\frac{-1}{p}\right) = (-1)^{\frac{p-1}{2}}$. In addition, $x^2 + 1 \equiv 0 \pmod{p}$ has two and only two solutions if it is solvable. This completes the proof. \square

Proposition 6. Let $p \in \mathbb{Z}$ be a prime and $l \in \mathbb{Z}, l \geq 1$. Then

$$v_3(p^l) = \begin{cases} 0 & \text{if } p \equiv 2 \pmod{3} \text{ or } 9|p^l, \\ 1 & \text{if } p = 3, \\ 2 & \text{if } p \equiv 1 \pmod{3}. \end{cases}$$

Proof. Let $(1, d) \in E_3(p^l)$ then $d^2 - d + 1 \equiv 0 \pmod{p^l}$. Since the system of two equations $x^2 - x + 1 \equiv 0 \pmod{p}$ and $2x - 1 \equiv 0 \pmod{p}$ has a common solution if $p = 3$, the number of solutions of $x^2 - x + 1 \equiv 0 \pmod{p^l}$ is equal to that of $x^2 - x + 1 \equiv 0 \pmod{p}$ if $p \neq 3$. The cases of $p = 2, 3$ or $9|p^l$ are trivial and we then let $p \geq 5$ in the following. $x^2 - x + 1 \equiv 0 \pmod{p}$ has a solution if $y^2 + 3 \equiv 0 \pmod{p}$ has a solution by taking $x = \frac{y+1}{2}$ and substituting $p - y$ for y when $y \equiv 0 \pmod{2}$. Then, $x^2 - x + 1 \equiv 0 \pmod{p}$

has a solution if $\left(\frac{-3}{p}\right) = 1$ if $p \equiv 1 \pmod{3}$ by

$$\left(\frac{-3}{p}\right) = \left(\frac{3}{p}\right) \left(\frac{-1}{p}\right), \left(\frac{3}{p}\right) = (-1)^{\frac{p-1}{2}} \left(\frac{p}{3}\right), \left(\frac{-1}{p}\right) = (-1)^{\frac{p-1}{2}}$$

and $\left(\frac{-3}{p}\right) = \left(\frac{p}{3}\right)$. In addition, $x^2 - x + 1 \equiv 0 \pmod{p}$ has two and only two solutions if it is solvable. This completes the proof. \square

The following results are well-known, see Proposition 1.43 in [1]. However, our proof is elementary and constructive.

Corollary 2. (1) $v_2(N) = \begin{cases} 0 & \text{if } 4|N, \\ \prod_{p|N} \left(1 + \left(\frac{-1}{p}\right)\right) & \text{otherwise.} \end{cases}$

(2) $v_3(N) = \begin{cases} 0 & \text{if } 4|N, \\ \prod_{p|N} \left(1 + \left(\frac{-3}{p}\right)\right) & \text{otherwise.} \end{cases}$

Proof. It is immediately from Theorem 4, Propositions 5 and 6. \square

Corollary 3. Let $g(N)$ be the genus of the modular curve $X_0(N)$. Then, for any $(N_1, N_2) = 1$,

$$g(N_1N_2) = 1 + \frac{\mu(N_1)\mu(N_2)}{12} - \frac{v_2(N_1)v_2(N_2)}{4} - \frac{v_3(N_1)v_3(N_2)}{3} - \frac{v_\infty(N_1)v_\infty(N_2)}{2}.$$

Proof. It is immediately from Theorems 1–3 and the formula for the genus of $X_0(N)$

$$g(N) = 1 + \frac{\mu(N)}{12} - \frac{v_2(N)}{4} - \frac{v_3(N)}{3} - \frac{v_\infty(N)}{2}.$$

\square

$E_3(N)$ can be constructed using Algorithm 3.

Algorithm 3: $E_3(N)$

(1) Construct $E_3(p^l)$ by general method; (2) Let $N = N_1N_2$ for $(N_1, N_2) = 1$. Given $E_3(N_1)$ and $E_3(N_2)$. $E_3(N)$ is constructed as follows. For all $(1, d_1) \in E_3(N_1)$, $(1, d_2) \in E_3(N_2)$, Determine d such that

$$d \equiv d_1 \pmod{N_1}, d \equiv d_2 \pmod{N_2} \text{ and } 1 \leq d \leq N.$$

Then, $(1, d) \in E_3(N)$ and all elements in $E_3(N)$ are constructed if all pairs in $E_3(N_1) \times E_3(N_2)$ are processed.

5. Concluding Remarks

In [7], Stein mentioned that another approach to list $\mathbb{P}^1(\mathbb{Z}/N\mathbb{Z})$ is to use that

$$\mathbb{P}^1(\mathbb{Z}/N\mathbb{Z}) \cong \prod_{p|N} \mathbb{P}^1(\mathbb{Z}/p^{v_p}\mathbb{Z}),$$

where $v_p = \text{ord}_p(N)$, and that it is relatively easy to enumerate the elements of $\mathbb{P}^1(\mathbb{Z}/p^n\mathbb{Z})$ for a prime power p^n . However, this approach had never been implemented by anyone as far as I know. Thus, Algorithm 1 in this paper could be regarded as an explicit implementation of Stein's ideas. All the algorithms described in this paper have been implemented in Wolfram Language, for these Wolfram programs, see [8]. We plan to rewrite these programs in the free open-source computer algebra system SAGE and incorporate them into Stein's program [9] or Walker's program [10].

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