



# Article Solution to a Conjecture on the Permanental Sum <sup>+</sup>

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**Abstract:** Let *G* be a graph with *n* vertices and *m* edges. A(G) and *I* denote, respectively, the adjacency matrix of *G* and an *n* by *n* identity matrix. For a graph *G*, the permanent of matrix (I + A(G)) is called the permanental sum of *G*. In this paper, we give a relation between the Hosoya index and the permanental sum of *G*. This implies that the computational complexity of the permanental sum is *NP*-complete. Furthermore, we characterize the graphs with the minimum permanental sum among all graphs of *n* vertices and *m* edges, where  $n + 3 \le m \le 2n - 3$ .

Keywords: permanental sum; Hosoya index; graph operation; extremal graph

MSC: 15A15; 05C31; 05E16

### 1. Introduction

The *permanent* of an *n* square matrix  $M = [m_{ij}]$  with  $i, j \in \{1, 2, ..., n\}$  is defined as

$$\operatorname{per}(M) = \sum_{\sigma \in \Lambda_n} \prod_{i=1}^n m_{i\sigma(i)},$$

where  $\Lambda_n$  denotes a symmetry group of order *n*. The computational complexity of the permanent of a matrix is #P-complete [1].

Let *G* be a graph on *n* vertices and *m* edges. Let A(G) and *I* denote the adjacency matrix of *G* and an *n* by *n* identity matrix, respectively. The *permanental polynomial* of *G* is defined as

$$\pi(G, x) = \operatorname{per}(xI - A(G)) = \sum_{k=0}^{n} b_k(G) x^{n-k},$$

where  $b_k(G)$  denotes the *k*th coefficient of  $\pi(G, x)$ . In particular,  $b_0(G) = 1$ . The permanental polynomials of graphs were first introduced in mathematics [2] and chemistry [3]. For more and additional information, see [4,5] and the references therein.

A Sachs graph *G* is a graph wherein each component is a single edge or a cycle. For a given an integer  $k \ge 0$ ,  $S_k(G)$  is the collection of all Sachs subgraphs  $H_k$  of order k in *G*, and let c(H) be the number of cycles in *H*. Merris et al. [2] gave a Sachs type result closely related to the coefficients of the permanental polynomial of *G* as below:

$$b_k(G) = (-1)^k \sum_{H \in S_k(G)} 2^{c(H)}, \ 0 \le k \le n.$$
<sup>(1)</sup>



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**Copyright:** © 2024 by the authors. Licensee MDPI, Basel, Switzerland. This article is an open access article distributed under the terms and conditions of the Creative Commons Attribution (CC BY) license (https:// creativecommons.org/licenses/by/ 4.0/). The *permanental sum* of *G*, written as PS(G), is the sum of the absolute values of all coefficients of  $\pi(G, x)$ , i.e.,

$$PS(G) = \sum_{k=0}^{n} |b_k(G)| = \sum_{k=0}^{n} \sum_{H \in S_k(G)} 2^{c(H)}.$$
(2)

Specifically, PS(G) = 1 if *G* is an empty graph. Wu and So [6] give an explicit formula for permanental sums as follows:

$$PS(G) = \operatorname{per}(I + A(G)),$$

which implies that calculating the permanental sum is #P-complete. The permanental sum of a graph was first considered by Tong [7]. In [8], Xie et al. captured a labile fullerene  $C_{50}(D_{5h})$ . Tong computed all 271 fullerenes in  $C_{50}$ . In his study, Tong found that the permanental sum of  $C_{50}(D_{5h})$  achieves the minimum among all 271 fullerenes in  $C_{50}$ . He pointed out that the permanental sum will be closely related to the stability of molecular graphs. For more information about permanental sums, see [9–12].

A *k*-matching in *G* is a set of *k* independent edges, and the number of *k*-matching is denoted by m(G, k). The matching number of *G*, denoted by v(G), is the maximum size of a matching in *G*. The *Hosoya index* Z(G) is defined as the total number of matchings of *G*, i.e.,

$$Z(G) = \sum_{k=0}^{\nu(G)} m(G,k).$$

In 1971, the chemist Hosoya firstly introduced Z(G) as a chemical application to describe the thermodynamic properties of saturated hydrocarbons. Later, the computational complexity of Z(G) was proved to be NP-complete [13,14]. Next, we introduce a relationship between the Hosoya index and the permanent below.

**Theorem 1.** Let *G* be a graph with *n* vertices, and let  $\mathscr{C}$  be the collection for which elements H' are disjoint unions of cycles in *G*. Then

$$PS(G) = Z(G) + \sum_{H' \in \mathscr{C}} 2^{c(H')} Z(G - H').$$
 (3)

**Proof.** By (1), we have

$$|b_k(G)| = \sum_{H \in S_k(G)} 2^{c(H)}, \ 0 \le k \le n.$$

$$\sum_{0}^{n} b_k^1(G) = Z(G).$$

Assume that H' is disjoint cycles with i vertices, and  $H' \in \mathscr{C}$ . For an integer k, the number  $H_k^2$  is equal to m(G - H', k - i). Thus,

$$b_k^2(G) = \sum_{H' \in \mathscr{C}} 2^{c(H')} m(G - H', k - i).$$

Combining the arguments as above, we have

$$PS(G) = \sum_{k=0}^{n} |b_k(G)| = \sum_{k=0}^{n} (b_k^1(G) + b_k^2(G)) = Z(G) + \sum_{H' \in \mathscr{C}} 2^{c(H')} Z(G - H').$$

By Theorem 1 and Jerrum's result [13], we have:

**Corollary 1.** The computation of PS(G) is NP-Complete.

Recently, some results on permanental sums were published. Let  $\mathscr{G}_{n,m}$  be a collection of connected simple graphs of order n and size m. Furthermore, if graph  $G \in \mathscr{G}_{n,m}$ and m = n + i, then G is called a tree, unicyclic graph, bicyclic graph, tricyclic graph, tetracyclic graph,  $\cdots$ , k-cyclic graph, where  $i = -1, 1, 2, 3, \cdots, n-3$ . In particular, every k-cyclic graph contains at least k cycles. Wu and Lai [15] determined the smaller bound of permanental sums of all unicyclic graphs in  $\mathscr{G}_{n,n}$ . And the corresponding extremal graphs  $F_n^n$  were determined, where the graph of  $F_n^n$  can be seen in Figure 1. Wu and Das [16] determined the lower bound of permanental sums of all bicyclic graphs in  $\mathscr{G}_{n,n+1}$ . And the corresponding extremal graphs  $F_{n+1}^n$  were determined, where graph of  $F_{n+1}^n$  can be seen in Figure 1. So et al. [6] characterized the lower bounds of the permanental sums of all tricyclic graphs in  $\mathscr{G}_{n,n+2}$ . And the corresponding extremal graphs  $F_{n+2}^n$  were determined, where the graph of  $F_{n+1}^n$  can be seen in Figure 1. According to the above results, So et al. [6] proposed a conjecture as follows.

**Conjecture 1.** For  $G \in \mathscr{G}_{n,m}$  with  $n + 3 \le m \le 2n - 3$ ,  $PS(G) \ge PS(F_n^m)$ , and the equality holds if and only if  $G = F_n^m$ . The graph of  $F_n^m$  can be seen in Figure 1.



**Figure 1.** Graph of  $F_n^m$ .

In this paper, we focus on Conjecture 1, and we give a solution as follows.

**Theorem 2.** Let  $G \in \mathscr{G}_{n,m}$ . If  $n + 3 \le m \le 2n - 3$ , then  $PS(G) \ge PS(F_n^m)$ , where the equality holds if and only if  $G \cong F_n^m$ .

The rest of this paper is organized as follows. In Section 2, we present some definitions and properties of permanental sums of graphs. In Section 3, we first present some graph operations that can be considered to be graph transformations, and show we that the transformed graph, generally, will have a smaller permanental sum than the original graph. Furthermore, we give the proof of Theorem 2. In the final section, we give a summary of this paper, and some new problems regarding permanent sums are introduced.

## 2. Preliminaries

All graphs considered in this work are undirected, finite, and simple graphs. For notation and terminology not defined here, see [17].

Let *G* be a graph with vertex set V(G) and edge set E(G). The *degree* of a vertex  $v \in V(G)$  is denoted by d(v). The *neighborhood* of vertex  $v \in V(G)$ , denoted by  $N_G(v)$ , is the

set of vertices adjacent to v. For a subgraph H of G, let G - V(H) (respectively, G - E(H)) denote the subgraph obtained from G by deleting the vertices and edges (respectively, deleting the edges) of H. In particular, if H is a vertex v (or an edge e), then G - V(H) (or G - E(H)) is written as G - v (or G - e). The path, cycle, and star of order n are denoted by  $P_n$ ,  $C_n$ , and  $S_n$ , respectively. Let  $G \cup H$  denote the union of two vertex disjoint graphs G and H. For any positive integer l, lG denotes the union of l disjoint copies of G.

Now we present some properties of permanental sums of graphs.

**Lemma 1** ([15]). *The permanental sum of a graph satisfies the following identities: (i) Let G and H be two graphs. Then* 

$$PS(G \cup H) = PS(G)PS(H).$$

(ii) Let uv be an edge of graph G, and C(uv) is the set of cycles containing uv. Then

$$PS(G) = PS(G - uv) + PS(G - v - u) + 2\sum_{C \in \mathcal{C}(uv)} PS(G - V(C)).$$

(iii) Let v be a vertex of graph G,  $N_G(v)$  be the set of neighbors of v, and C(v) be the set of cycles containing v. Then

$$PS(G) = PS(G-v) + \sum_{u \in N_G(v)} PS(G-v-u) + 2\sum_{C \in \mathcal{C}(v)} PS(G-V(C)).$$

**Lemma 2** ([18]). Let G be a connected simple graph with n vertices. Then  $n \leq PS(G) \leq n!$ . The left equality holds if and only if  $G \cong S_n$ , and the right equality holds if and only if  $G \cong K_n$ .

**Lemma 3** ([6]). Let  $G \in \mathscr{G}_{n,n+2}$ . Then  $PS(G) \ge PS(F_n^{n+2})$ , where the equality holds if and only if  $G \cong F_n^{n+2}$ .

**Lemma 4** ([19]). Suppose that  $G \in \mathcal{G}_{n,m}$ , where  $n + 2 \le m \le 2n - 3$ , and all cut edges are pendent edges incident with the same vertex. Then there exists an edge uv in G such that the subgraphs G - uv and  $G - \{u, v\}$  are still connected.

### 3. Proof of Theorem 2

Before we prove Theorem 2, we introduce three graph operations that can be considered to be graph transformations, and we show that, generally, the transformed graph will have smaller permanental sum than the original graph.

**Definition 1.** Let u be a vertex of graph  $G_0$ . Denote by  $G_1$  the graph obtained from  $G_0$  and a tree T by attaching u to a vertex v of T. Denote by  $G_2$  the graph obtained from  $G_0$  and a star  $S_{|T|}$  by attaching u to the center v' of  $S_{|T|}$ . We designate the transformation from  $G_1$  to  $G_2$  as type I.

**Lemma 5** ([15]). Suppose that  $G_1$  and  $G_2$  are two graphs as defined by Definition 1. Then  $PS(G_1) \ge PS(G_2)$ , and the equality holds if and only if T is a star and v is the center of T.

**Definition 2.** Let  $G_0$  be a graph of order of at least 2, and let  $u, v \in G_0$ . Denote by G(s, t) the graph obtained from  $G_0$  by attaching  $s \ge 1$  and  $t \ge 1$  pendant vertices to u and v, respectively. Denote by G'(s+t) the graph obtained from  $G_0$  by attaching s + t pendant vertices to u, and denote by G''(s+t) the graph obtained from  $G_0$  by attaching s + t pendant vertices to v. The resulting graphs G(s,t), G'(s+t), and G''(s+t) are displayed in Figure 2. We designate the transformation from G(s,t) to G'(s+t) or G''(s+t) as type II.

**Lemma 6.** Suppose that G(s,t), G'(s+t), and G''(s+t) are three graphs defined as in Definition 2. Then PS(G(s,t)) > PS(G'(s+t)) or PS(G(s,t)) > PS(G''(s+t)). **Proof.** By Lemma 1, deleting  $e_1, ..., e_s, h_1, ..., h_t$  one by one in G(s, t), we get that

$$\begin{array}{ll} PS(G(s,t)) \\ = & PS(G(s,t) - e_1) + PS(G(s,t) - u_1 - u) \\ = & PS(K_1)PS(G(s-1,t)) + PS((s-1)K_1)PS(G(0,t) - u) \\ = & PS(G(s-1,t) - e_2) + PS(G(s-1,t) - u_2 - u) + PS(G(0,t) - u) \\ = & PS(G(s-2,t)) + PS(G(0,t) - u) + PS(G(0,t) - u) + PS(G(0,t) - u) \\ = & PS(G(s-2,t)) + PS(G(0,t) - u) + PS(G(0,t) - u) \\ = & PS(G(s-2,t)) + 2PS(G(0,t) - u) \\ = & \dots \\ = & PS(G(s-(s-1),t) - e_s) + PS(G(s-(s-1),t) - u_s - u) + (s-1)PS(G(0,t) - u) \\ = & PS(G(s,t)) + 2PS(G(0,t) - u) + (s-1)PS(G(0,t) - u) \\ = & PS(G(0,t)) + SPS(G(0,t) - u) + (s-1)PS(G(0,t) - u) \\ = & PS(G(0,t)) + SPS(G(0,t) - u) + (s-1)PS(G(0,t) - u) \\ = & PS(G(0,t-1)) + PS(G(0,t) - v_1 - v) + SPS(G(0,t) - u) \\ = & PS(G(0,t-1)) + PS(G(0,t) - v_1 - v) + SPS(G(0,t) - u) \\ = & PS(G(0,t-1)) + PS(G(0,t-1) - v_2 - v) + PS(G(0,v) + sPS(G(0,t) - u) \\ = & PS(G(0,t-1)) + PS(G(0,t-1) - v_2 - v) + PS(G_0 - v) + sPS(G(0,t) - u) \\ = & PS(G(0,t-2)) + 2PS(G(0,t-1) - v_2 - v) + PS(G_0 - v) + sPS(G(0,t) - u) \\ = & PS(G(0,t-2)) + 2PS(G_0 - v) + sPS(G(0,t) - u) \\ = & \dots \\ = & PS(G(0,t-2)) + 2PS(G_0 - v) + sPS(G(0,t) - u) \\ = & PS(G(0,t-2)) + 2PS(G_0 - v) + sPS(G(0,t) - u) \\ = & PS(G(0,t-2)) + 2PS(G_0 - v) + sPS(G(0,t) - u) \\ = & PS(G(0,t-2)) + 2PS(G_0 - v) + sPS(G(0,t) - u) \\ = & PS(G(0,t-2)) + 2PS(G_0 - v) + sPS(G(0,t) - u) \\ = & PS(G(0,t-2)) + 2PS(G_0 - v) + sPS(G(0,t) - u) \\ = & PS(G(0,t-2)) + PS(G_0 - v) + sPS(G(0,t-1) - u) + PS(G(0,t) - u - v_1 - v)] \\ = & PS(G_0) + tPS(G_0 - v) + s[PS(G(0,t-1) - u_1) + PS(G(0,t) - u - v_1 - v)] \\ = & PS(G_0) + tPS(G_0 - v) + s[PS(G(0,t-1) - u - h_2) + PS(G(0,t-1) - u - v_2 - v) + PS(G_0 - u - v)] \\ = & PS(G_0) + tPS(G_0 - v) + s[PS(G(0,t-1) - u - h_2) + PS(G(0,t-1) - u - v_2 - v) + PS(G_0 - u - v)] \\ = & PS(G_0) + tPS(G_0 - v) + s[PS(G(0,t-2) - u) + 2PS(G_0 - u - v)] \\ = & PS(G_0) + tPS(G_0 - v) + s[PS(G(0,t-1) - u - h_1) + PS(G(0,t-1) - u - v_2 - v) + PS(G_0 - u - v)] \\ = & PS(G_0) + tPS(G_0 - v) + s[PS(G_0 - u + v)] \\ = & PS(G_0) + tPS(G_0 - v) + s[PS(G_0 - u + v)] \\ = & PS(G_0) + tPS(G_0 - v) + s[PS(G_0 - u + v)] \\ = & PS(G_0) + tPS(G_0 - v)$$

Similarly, by Lemma 1, deleting  $e_1, ..., e_s, h_1, ..., h_t$  one by one in G'(s + t), we obtain that

$$\begin{split} PS(G'(s+t)) \\ &= PS(G'(s+t) - e_1) + PS(G'(s+t) - u_1 - u) \\ &= PS(K_1)PS(G'(s+t-1)) + PS((s+t-1)K_1)PS(G_0 - u) \\ &= PS(G'(s+t-1)) + PS(G_0 - u) \\ &= PS(G'(s+t-1) - e_2) + PS(G'(s+t-1) - u_2 - u) + PS(G_0 - u) \\ &= PS(K_1)PS(G'(s+t-2)) + PS((s+t-2)K_1)PS(G_0 - u) + PS(G_0 - u) \\ &= PS(G'(s+t-2)) + PS(G_0 - u) + PS(G_0 - u) \\ &= PS(G'(s+t-2)) + 2PS(G_0 - u) \\ &= \dots \\ &= PS(G'(s+t-(s+t-1) - h_t) + PS(G'(s+t-(s+t-1) - v_t - u) \\ &\quad + (s+t-1)PS(G_0 - u) \\ &= PS(K_1)PS(G_0) + PS(G_0 - u) + (s+t-1)PS(G_0 - u) \\ &= PS(G_0) + PS(G_0 - u) + (s+t-1)PS(G_0 - u) \\ &= PS(G_0) + (s+t)PS(G_0 - u). \end{split}$$

By the symmetry of the calculation of PS(G'(s + t)), it is easy to obtain that

$$PS(G''(s+t)) = PS(G_0) + (s+t)PS(G_0 - v).$$

Direct calculation yields

$$\Delta_1 = PS(G(s,t)) - PS(G'(s+t)) = t[PS(G_0 - v) - PS(G_0 - u) + sPS(G_0 - u - v)], \Delta_2 = PS(G(s,t)) - PS(G''(s+t)) = s[PS(G_0 - u) - PS(G_0 - v) + tPS(G_0 - u - v)].$$

If  $\triangle_1 \leq 0$ , then  $PS(G(s,t)) \leq PS(G'(s+t))$ , and so  $PS(G_0-u) \geq PS(G_0-v) + sPS(G_0-u-v)$ . Thus,

$$\Delta_2 = PS(G(s,t)) - PS(G''(s+t)) \geq s[PS(G_0 - v) + sPS(G_0 - u - v) - PS(G_0 - v) + tPS(G_0 - u - v)] = s(s+t)PS(G_0 - u - v) > 0.$$

If  $\triangle_2 \leq 0$ ,  $PS(G(s,t)) \leq PS(G''(s+t))$ , and so  $PS(G_0 - v) \geq PS(G_0 - u) + tPS(G_0 - u - v)$ . Hence,

$$\Delta_1 = PS(G(s,t)) - PS(G'(s+t)) \geq t[PS(G_0 - u) + tPS(G_0 - u - v) - PS(G_0 - u) + sPS(G_0 - u - v)] = t(s+t)PS(G_0 - u - v) > 0.$$

So PS(G(s,t)) > PS(G'(s+t)) or PS(G(s,t)) > PS(G''(s+t)).  $\Box$ 



**Figure 2.** Graphs G(s,t), G'(s+t), and G''(s+t).

**Definition 3.** Let G and H be two disjoint connected graphs of order of at least 2 with  $v \in V(G)$ and  $u \in V(H)$ . Denote by  $G_1$  the graph obtained from the union of G and H by adding a new edge uv. Let  $G_2$  denote the graph obtained from  $G_1$  by deleting the edge uv and identifying u with v to form a new vertex x and attaching a pendent vertex y to x. The resulting graphs  $G_1$  and  $G_2$  are displayed in Figure 3. We designate the transformation from  $G_1$  to  $G_2$  as type III.



**Figure 3.** Graphs  $G_1$  and  $G_2$ .

**Lemma 7.** Let  $G_1$  and  $G_2$  be two graphs as defined by Definition 3. Then  $PS(G_1) > PS(G_2)$ .

**Proof.** By Lemma 1, we have

$$\begin{aligned} &PS(G_{1}) \\ &= PS(G)PS(H) + PS(G-v)PS(H-u) \\ &= 2PS(G-v)PS(H-u) + PS(G-v)\sum_{u'\in N_{H}(u)} PS(H-\{u,u'\}) + 2PS(G-v) \\ &\times \sum_{C\in\mathcal{C}_{H}(u)} PS(H-V(C)) + \sum_{v'\in N_{G}(v)} PS(G-\{v,v'\})PS(H-u) \\ &+ \sum_{v'\in N_{G}(v)} PS(G-\{v,v'\})\sum_{u'\in N_{H}(u)} PS(H-\{u,u'\}) + 2\sum_{v'\in N_{G}(v)} PS(G-\{v,v'\}) \\ &\times \sum_{C\in\mathcal{C}_{H}(u)} PS(H-V(C)) + 2\sum_{C\in\mathcal{C}_{G}(v)} PS(G-V(C))PS(H-u) \\ &+ 2\sum_{C\in\mathcal{C}_{G}(v)} PS(G-V(C))\sum_{u'\in N_{H}(u)} PS(H-\{u,u'\}) + 4\sum_{C\in\mathcal{C}_{G}(v)} PS(G-V(C)) \\ &\sum_{C\in\mathcal{C}_{H}(u)} PS(H-V(C)). \end{aligned}$$

Similarly, by Definition 3,  $G_2 - x = (G - v) \cup (H - u)$ ,  $N_{G_2}(x) = N_G(v) \cup N_H(u) \cup \{y\}$ , and by Lemma 1, we get that

$$\begin{split} & PS(G_2) \\ &= PS(G_2 - xy) + PS(G_2 - \{x, y\}) \\ &= PS(G_2 - x) + \sum_{x' \in N_{G_2}(x)/y} PS(G_2 - \{x, x'\}) + 2\sum_{C \in \mathcal{C}_{G_2}(x)} PS(G_2 - V(C)) \\ &+ PS(G_2 - \{x, y\}) \\ &= 2PS(G - v)PS(H - u) + \sum_{x' \in N_G(v)} PS(G - \{x, x'\}) + \sum_{x' \in N_H(u)} PS(H - \{x, x'\}) \\ &+ 2\sum_{C \in \mathcal{C}_G(v)} PS(G - V(C)) + 2\sum_{C \in \mathcal{C}_H(u)} PS(H - V(C)) \\ &= 2PS(G - v)PS(H - u) + PS(H - u) \sum_{v' \in N_{G(v)}} PS(G - \{v, v'\}) + PS(G - v) \\ &\times \sum_{u' \in N_H(u)} PS(H - \{u, u'\}) + 2\sum_{C \in \mathcal{C}_G(v)} PS(G - V(C)) + 2\sum_{C \in \mathcal{C}_H(u)} PS(H - V(C)). \end{split}$$

Thus

$$PS(G_{1}) - PS(G_{2}) = 2 \sum_{C \in \mathcal{C}_{H}(u)} PS(H - V(C))(PS(G - v) - 1) + 2 \sum_{C \in \mathcal{C}_{G}(v)} PS(G - V(C)) \\ \times (PS(H - u) - 1) + \left[ \sum_{v' \in N_{G}(v)} PS(G - \{v, v'\}) + 2 \sum_{C \in \mathcal{C}_{G}(v)} PS(G - V(C)) \right] \\ \times \left[ \sum_{u' \in N_{H}(u)} PS(H - \{u, u'\}) + 2 \sum_{C \in \mathcal{C}_{H}(u)} PS(H - V(C)) \right] \\ > 0.$$

**Remark 1.** For topological indices of graphs, they have similar graph operations as above. For example, the Hosoya index [19], Wiener index [20,21], etc.

**Proof of Theorem 2.** Suppose that  $G \in \mathcal{G}_{n,m}$  when  $n + 3 \le m \le 2n - 3$ . By repeatedly applying the transformations *I*, *II*, and *III* and by Lemmas 5–7, we can obtain a graph G' from *G* such that all bridges are pendent edges incident with the same vertex and  $PS(G') \le PS(G)$ , where the equality holds if and only if  $G' \cong G$ . Additionally, by Lemma 4, there exists an edge uv in G' such that both  $G' - uv \in \mathcal{G}_{n,m-1}$  and  $G' - \{u, v\}$  are connected.

We use induction on *m*. Assume that m = n + 3, i.e., G' is a tetracyclic graph. On the one hand,

$$PS(F_n^{n+3}) = PS(F_n^{n+3} - uv) + PS(F_n^{n+3} - \{u, v\}) + 2\sum_{C \in \mathcal{C}(uv)} PS(F_n^{n+3} - V(C))$$
  
=  $PS(F_n^{n+2}) + PS(S_{n-2}) + 2 \times 4.$ 

On the other hand, since G' contains at least 4 cycles,  $G' - uv \in \mathcal{G}_{n,n+2}$ , and  $G' - \{u, v\}$  is a connected graph of order n - 2, then

$$PS(G) \geq PS(G') \\ = PS(G' - uv) + PS(G' - \{u, v\}) + 2\sum_{C \in \mathcal{C}(uv)} PS(G' - V(C)) \\ \geq PS(F_n^{n+2}) + PS(S_{n-2}) + 8 \\ = PS(F_n^{n+3}),$$

where the last inequality is derived from Lemmas 2 and 3. In order for the equalities to hold, all the inequalities above should be equalities. Then  $G' - uv \cong F_n^{n+2}, G' - \{u, v\} \cong S_{n-2}, \sum_{C \in \mathcal{C}(uv)} PS(G' - V(C)) = 4$  and  $G \cong G'$ . So it is not hard to verify that  $G \cong G' \cong F_n^{n+3}$ .

Suppose now that the statement holds for m - 1 ( $m \ge n + 4$ ). We will prove this for m as follows. We have

$$PS(F_n^m) = PS(F_n^m - uv) + PS(F_n^m - \{u, v\}) + 2\sum_{C \in \mathcal{C}(uv)} PS(F_n^m - V(C))$$
  
=  $PS(F_n^{m-1}) + PS(S_{n-2}) + 2(m - n + 1).$ 

Note that G' consists of at least m - n + 1 cycles. It follows that  $G' - uv \in \mathscr{G}'_{n,m-1}$ , and  $G' - \{u, v\}$  is a connected graph of order n - 2.

$$PS(G) \geq PS(G') \\ = PS(G' - uv) + PS(G' - \{u, v\}) + 2\sum_{C \in \mathcal{C}(uv)} PS(G' - V(C)) \\ \geq PS(F_n^{m-1}) + PS(S_{n-2}) + 2(m - n + 1) \\ = PS(F_n^m),$$

where the last inequality is derived from the induction hypothesis and Lemma 2. Similarly, to make the equalities hold, all of the inequalities above should be equalities. Then  $G' - uv \cong F_n^{m-1}$ ,  $G' - \{u, v\} \cong S_{n-2}$ ,  $\sum_{C \in \mathcal{C}(uv)} PS(G' - V(C)) = m - n + 1$  and  $G \cong G'$ . Therefore, it is not hard to verify that  $G \cong G' \cong F_n^m$ . Hence, the assertion holds for *m*. Consequently,

It is not hard to verify that  $G \cong G \cong F_n^m$ . Hence, the assertion holds for *m*. Consequently it holds for all  $n + 3 \le m \le 2n - 3$ .

This completes the proof.  $\Box$ 

#### 4. Summary

In this paper, we prove that the computational complexity of a permanental sum is *NP*-complete. In particular, we determine the minimum value of permanental sums of all graphs with given *n* vertices and  $n + 3 \le m \le 2n - 3$  edges. This result promotes the study of permanental sums. It raises a lot of interesting questions, such as those related to determining the sharp bound of a permanental sum of all graphs in  $\mathcal{G}_{n,m}$  if  $n + 3 \le m \le 2n - 3$  and questions regarding characterizing the bound of permanental sums of all graphs in  $\mathcal{G}_{n,m}$  if m > 2n - 3, etc.

A permanent is a generalized matrix function that has important applications in chemistry [22–24]. A permanental sum is a derivative of a permanent, and it is a topological index proposed to explain special chemical phenomena from a mathematical point of view. The result in Theorem 2 is interesting for chemistry. In [25], the authors pointed out that every graph with a maximum degree that is no more than 4 has a chemical molecule corresponding to it. The result in Theorem 2 implies that the smaller bound of permanental sums of all chemical molecules is determined. And an interesting problem arises, i.e., characterizing the sharp bound of permanental sums of chemical molecules.

In conclusion, the above problems will guide us to continue our research.

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